

TRANSFER OF AN INTEGER BINARY VALUE INTO DECIMAL ASCII STRING

Abstract

The objective of this article is to present the variety of methods available for transfer of an integer binary number into the decimal ASCII string. The common methods often are not the fastest ones and there are some almost forgotten and still reliable methods (by repetitive subtractions and additions) and some less known ones like division via multiplication or binary tree search for the decimal digits.

Although we write from left to right, the weight places in decimal numbers are organized in reverse order, from right to left, and this requires minor modifications and additional code in some algorithms. It is much more convenient on paper or papyrus to write from left to right because the process of writing can be visually controlled and the angle of the pen is higher than 90° which facilitates pen's gliding on paper, while writing on the clay has to be performed in reverse order because the angle of the pen, i.e. wedge shaped stick is less than 90° and therefore it is in much better position to scratch and punch the raw clay table. The reverse order of digits in decimal numbers is the legacy of that ancient era of cuneiform writing on clay tables which became quite obvious during the development of these routines for 2→10 radix conversion. This challenges historical theory that decimal numbers appear firstly in India because they wrote from left to right on fabrics and frescoes and therefore decimal numbers could be much older and with a different origin.

The problem itself is mathematically quite clear and essentially it consists of radix transition from radix 2 into the radix 10, but its practical solution is not that simple.

There are 6 basic ways to transfer register's binary value into its decimal ASCII string:

1. Successive subtraction of decimal weight coefficients (1, 2, 4, 8, 16, 32, 64, 128, 256, etc.) and countering till the result is being negative: This method yields the decimal digits in direct order with leading zeroes that should be truncated, but the number of these leading zeroes can be vastly reduced by an estimation of the leading zero positions with the error of eventually a single position only. The set of weight coefficients could be stored in an array and then grabbed with **LODSB** instruction or on the stack and manipulated via **POP** one. The average number of iterations per decimal digit is 4.5 loops. The method is quite simple and yet it utilizes only fast instructions for integer addition and subtraction. This method was very popular in the era of 8-bit computers when it was almost only available option due to lack of complex assembly instructions. The routines based on that method are **EAX2SUB** and **ESI2ASC** ones. **EAX2SUB** pulls the weight coefficients from an array while **ESI2ASC** uses the stack. The array based routine is faster than the stack one because it can finish the task earlier, which is not the case with the stack based routine. The stack based **ESI2ASC** routine has to

pull all the coefficients from the stack to be able to return back with **RET** instruction while the array based **EAX2SUB** does not have such limitation and it can cancel its main loop at any time. The array based routine is also suitable for one more optimization: the number of decimal digits can be estimated in regards to the most significant bit of the converted binary number which is determined by the **BSR** instruction. Thereby the **NUMDIG** array contains the number of certainly leading zeroes in respect to the most significant bit of the converted binary number and that estimation, as aforesaid, can miss for only one leading zero digit's position which utterly eliminates usage of **REP SCASB** instruction's pair for truncation of leading zero's decimal digits, just as it is done in the **ESI2ASC** stack based routine. Disadvantage of the array based method is that it alters more registers than the stack's based one and it also uses **CMOVZ** instruction to clear the undefined state of destination register of **BSR** instruction when the source register is zero, and this instruction is not available in most embedded systems because it firstly appeared within Pentium Pro processor.

2. Addition of decimal numbers via **AAA** instruction for unpacked BCD arithmetic: The essence of this approach is consisted of the fact that we can convert the whole decimal string by adding single decimally digits of the binary positions (1, 2, 4, 8, 16, 32, etc) which can be done by BCD instructions. There is a handful **AAA** instruction that is utilized here and this is instruction for decimal correction after decimal digit per digit addition of decimally unpacked numbers. Decimally unpacked number is consisted of one decimal digit per a byte while decimally packed number is consisted of two decimal digits per byte where every single nibble of the byte holds a decimal digit. So, if we decimally add all numbers whose denote weight binary numbers (1, 2, 4, 8, 16, 32, 64, 128, etc) for proper binary positions, then we will obtain the decimal string of the whole binary input. Example: 25 is $25D = 11001B = 1D + 8D + 16D$, $1 + 8 = 9 + 0 \cdot 10$, $9 + 16 = 9 + 6 + 1 \cdot 10 = 5 + 1 \cdot 10 + 1 \cdot 10 = 25$. This method is not very fast but it is extremely suitable for VHDL hardware implementation. These software realizations are carefully coded and therefore the current implementations minimize the number of operation by omission of processing of leading zero positions which is done in an attempt to make this advance method more comparable by speed to other theoretically much simpler ones. The routines that demonstrate this method are **EDX2DEC** and **EAX2BCD** ones. **EDX2DEC** routine is entirely based on the **AAA** machine instruction and usage of this instruction can be avoided by utilization of an array which might be better solution because BCD arithmetic is obsolete in X64 mode and therefore **AAA** instruction does not exist in the 64-bit world anymore, and also, such array can instantly hold ASCII values of digits making the routine able to process ASCII digits directly. The routine can operate either by the successive shifting of the originating number one binary position right and then to check whether the carry flag is switched on, or by utilization of the **BT** instruction – there is chosen algorithm with shifting. This method also extracts digits in direct order

with eventually an extra leading zero which truncation is quite simple. **EAX2BCD** routine that operates without **AAA** instruction is derived from **EDX2DEC** one and it does not use obsolete **AAA** machine instruction. Instead of **AAA** instruction **EAX2BCD** routine utilizes the **ADDIT** array that holds all necessary decimal positions for addition of two single digit decimal numbers with the decimal carry. The **EAX2BCD** routine also deals directly with ASCII values of decimal digits because the weight array already contains ASCII digits making this routine faster than the **EDX2DEC** one based on **AAA** instruction.

3. The extraction via division by 10: This method is very convenient because **DIV** instruction simultaneously yields both the result and the remainder. Thereby the decimal digits are extracting in reverse order by repetitive division by 10 – the remainder is a single decimal digit and the result is something that is going to be used once again in the following iteration. The digits are extracting in reverse order and therefore they should be storing backward. The method is demonstrated by the **EAX2DEC** routine. The disadvantage of the **DIV** assembly instruction is that it alters Zero flag incorrectly and thus zero flag is not switched on whenever the result of the integer division is zero and therefore one additional **TEST** instruction must be used to check whether the result is zero because this zero result terminates the decimal digits' extraction.
4. The extraction via division by 10 by multiplication: A strange property of the integer numbers is that division by the constant can be achieved via multiplication with another appropriate constant. For the case of the 32-bit variable the multiplication constant is 858993459 ($x/10 \approx (858993459 \cdot (x+1)) \backslash 2^{33}$), for the case of 16-bit arithmetic the constant is 13107 ($x/10 \approx 13107 \cdot ((x+1)) \backslash 2^{17}$) and for 64-bit arithmetic the corresponding formula is $x/10 \approx (3689348814741910323 \cdot (x+1)) \backslash 2^{65}$. Disadvantage of this method is that there is no remainder denoting single decimal digits and therefore it has to be additionally calculated. The advantage is usage of **MUL** instruction which is much faster than **DIV** one and therefore the entire subroutine is much faster too. The routine that demonstrates this is **EAX2ASC** one. The improvement of the speed by the replacement of **DIV** instruction is huge although it may happen in the future that it will not be the case – the latency of the **DIV** instruction may reach the latency of **MUL** instruction and then this meandering approach will be obsolete making **EAX2ASC** routine even slower than **EAX2DEC** one. The extraction of the decimal digits is performing in reverse order and therefore the digits have to be storing reversely.
5. The decimal digits can be extracted from binary numbers by the Binary Tree Search method: The theory says that the number of comparisons in binary tree per digit should be equal to $\text{LOG}_2(10)$, and in our case the maximal length of a single branch is maximally 4 comparisons ($\text{LOG}_2(10) \approx 3.32$). This method can be implemented via double indirect addressing (routine **EAX2BTR**) or by utilization of pair of complementary **CMOVxx** instructions (**EAXCMOV**). The **CMOVxx** instructions on some processors have a serious flaw in which the source memory address will be read regardless the condition is fulfilled

or not and then the contents will be transferred to the destination location in respect to the condition. In this particular case it means that Access Violation Fault will be launched because the first **CMOVC** instruction (from the pair of complementary ones) will affect the register holding the indirect memory address of the second **CMOVNC** instruction which is not expected to be performed at all according the condition but it will be and that launches this fault – this moment is clearly commented in the **EAXCMOV** routine. Furthermore, the overflow fault may be also triggered in the situation where the source address is calculated by doubling or quadrupling of the addressing register already altered with wrong content. So, the pair of complementary **CMOVC/CMOVNC** instructions must be kept isolated with preserved address and that requires usage of an extra register. Although the reason for such persistent ignoring of this error is apparently inscrutable, it might be deliberately done because there is a perceivable reason for this failure – on this way the queue of instructions is not flushed and thereby the execution is faster. This routine is consisted of two stages: on the first stage the number of decimal digits is determining and on the second stage the extraction of these digits is performing. As it is mentioned in the above text, the name of the routine which utilizes the double indirect addressing is **EAX2BTR** and of the routine which uses **CMOVC/CMOVNC** instructions is **EAXCMOV**. The **LB** array contains the guiding rules for the search in binary tree while **D** array contains the digits weight coefficients used for subtractions during binary tree search. The method of Binary Tree Search requires substantial effort and dedication to be involved in its design and implementation and therefore the proper debugger and patient are necessary and DOS Box environment may help a lot in such situations.

6. Math coprocessor utilization: There is a method that is based on the utilization of the numeric coprocessor. Although the algorithm itself is hidden behind particular coprocessor instruction, this method still remains as a plausible option. The coprocessor contains one handful **FBSTP** instruction that yields packed BCD digits of the number. So, the integer should be converted into the double precision real number (i.e. to be put it into a coprocessor's register) and then to be transferred back from coprocessor as the packed BCD number. This is regularly doing by only two coprocessor's instructions: **FILD** and **FBSTP** respectively. The rest of the job is unpacking these packed decimal digits and to add 48 (Chr\$(48)="0") to each one, i.e. to allocate them all correct ASCII values. This process can be performed on two ways – via the array or by **SHRD** instruction. The **SHRD** instruction is able to shift pair of registers simultaneously for a nibble (denoting a packed decimal digit), i.e. 4 bits right. The disadvantage of the array utilization is that it simultaneously yields pair of digits while the disadvantage of the **SHRD** instruction usage is that this instruction is not able to set zero flag properly when both registers are null, which then requires utilization of an extra register. The routine that utilizes an array is the **EAX2AST** and the routine based on the **SHRD** instruction is **EAX2AFL**. The **FBSTP** instruction may be very slow depending of the particular processor (its

latency is 164 clocks on new processors) and the unpacking process also stands for while and therefore the usage of the **FBST** instruction should be carefully evaluated. The advantage of this method is that it is able to simply convert even 64-bit number. This extraction could be performed in both directions and reverse one is better solution when there is no prior knowledge of number of decimal digits that should be extracted. The table based routine also yields the leading zeroes which should be truncated either by additional **REPE SCASB** instructions pair or by an estimation of decimal digits as it is done in **EAX2SUB** routine, and the truncating of leading zeroes is significant disadvantage of this method. The number of digits can be exactly predetermined numerically as $1 + \text{INT}(\text{LOG}_{10}(2) \cdot \text{LOG}_2(n))$, which is usually doing by the following sequence of coprocessor instructions: **FLDLG2**, **FILD**, **FYL2X** and **FISTP** followed by **INC** integer instruction. The disadvantage of this approach is that it launches an overflow error for the zero input because the logarithm of the zero is infinite. This numerical computation is also slow and should be used only on bigger integers where the additional advantage could be achieved later in the main loop.

7. "Lazy" method: This is the method which is not the real one or at least less decent one because it utilizes existing Windows API function **Wsprintf** that is able to convert integer contents of a register into corresponding decimal ASCII value. The method is tremendously slow and it should be used only when the speed is less limiting factor than the time itself. The routine based on this API call is **EAX2WIN** one. The **Wsprintf** is full of various formatting and rich with options of all sorts which make this routine tremendously slow. This routine really raises the fear that various compilers may routinely use it tomorrow...

The demo program also contains the **SQR** routine which simultaneously yields both the result and the remainder. The explanation of the algorithm is beyond the scope of this article and therefore its explanation of operation will be omitted here, but let say that it is derived from the following formula $(x + \Delta x)^2 = x^2 + 2 \cdot x \cdot \Delta x + (\Delta x)^2$ and basically this is a kind of exhaustion of the quantified interval with Binary Tree Search.

The Windows assembly programming is very meticulous job because Windows environment is much more sensitive than the DOS one and therefore the direction flag must be maintained clear all the time otherwise some Window API call may fail or console print may fail, or some other exotic error may be triggered and for this reason another development environment may help a lot and DOS in DOS Box is close enough.

The another disadvantage of the Windows environment in respect to DOS one is the strong isolation between the **DATA** and **CODE** sections which impedes coding of the self modifying routine and also disables the possibility of keeping the routines and their memory variables together in the listing. Also, in an **OBJ** file usually rather exists **.text** section than the true **.CODE** section although existence of the **.CODE** section should be expected according the common myth available in most of Assembler's textbooks. The ability of keeping the routines gathered with their constants and variables is very important for the Cat & Paste rapid subroutines assembling. These

irrational limitation can be surpassed by the modification of the Text section properties via altering of the `.text` section's flags through the section's command of linker: `Link.exe /section:".text",ERW`, which also duly debunks the second myth frequently cited in Assembly textbooks that the self-modifying code is not possible within Windows and on newer processors.

Although this method alters section's read/write property via appropriate linker commands and therefore it should be avoided because its implementation requires certain reader's experience based on the particular assembler's peculiarities (ML, jwasm, etc.) that may distract reader from the essence of the presented algorithms, this particular way is chosen because it offers creation of much more readable listing. All these routines were initially written in DOS and then transferred in Windows environment via WinASm IDE with MASM32 10.0 package and ML 6.14.8444 with Link 5.12.8078. The alternative assembler used here was JWASM v2.04c.

The routines in this text are already coded highly optimized and the numbers of instructions are minimized. Any assembly routine is much faster than any higher language corresponding one just because the contemporary compilers do not utilize all the instructions and do not spread all data through the available registers. The MMX and XMM registers usually remain unused or just partially populated by the contemporary compilers and the code is exceedingly filled with `MOV` instructions, thereby carefully crafted assembly routines usually run much faster.

The listing of the aforesaid routines follows:

```
;Author of the program and of all algorithms is Andrija Radovic,
;All Rights Reserved, ©2010.
;This code should be assembled with:
;masm32\bin\ML.EXE /c /coff /Cp /nologo /I\Masm32\Include ASCIIwin.asm
;and linked with the following line:
;masm32\bin\LINK.EXE /SUBSYSTEM:CONSOLE /RELEASE /VERSION:4.0 /LIBPATH:"\Masm32\Lib" /section:".text",ERW
ASCIIwin.obj /OUT:ASCIIwin.exe

.686
.MODEL flat, stdcall
OPTION casemap:none ;case sensitive
INCLUDE \masm32\include\windows.inc
INCLUDE \masm32\include\kernel32.inc
INCLUDELIB \masm32\lib\kernel32.lib
INCLUDE \masm32\include\user32.inc
INCLUDELIB \masm32\lib\user32.lib
INCLUDE \masm32\include\masm32.inc
INCLUDELIB \masm32\lib\masm32.lib

;-----

.DATA
TESTS DB "This is the test...", 13, 10, 0
INTRES DB 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0
RESULT DB 0

;-----

.STACK 20000

;-----

.CODE
ASCII:
    INVOKE StdOut, ADDR TESTS
    CALL TESTR
    INVOKE StdOut, ADDR TESTS
    CALL MAINN
    INVOKE StdOut, ADDR TESTS
    INVOKE ExitProcess, 0

;-----

TESTV DD 0

TESTR PROC
    MOV TESTV, 3000000000
    MOV EDX, TESTV
    CALL EDX2DEC
    CALL NEW_ROW
    MOV EAX, TESTV
    ADD EAX, 1
    CALL EAX2SUB
    CALL NEW_ROW
    MOV ESI, TESTV
TESTR ENDP
```

```

        ADD     ESI, 2
        CALL    ESI2ASC
    CALL NEW_ROW
    MOV     EAX, TESTV
        ADD     EAX, 3
        CALL    EAX2ASC
    CALL NEW_ROW
    MOV     EAX, TESTV
        ADD     EAX, 4
        CALL    EAX2AST
    CALL NEW_ROW
    MOV     EAX, TESTV
        ADD     EAX, 5
        CALL    EAX2AFL
    CALL NEW_ROW
    MOV     EAX, TESTV
        ADD     EAX, 6
        CALL    EAX2DEC
    CALL NEW_ROW
    MOV     EAX, TESTV
        ADD     EAX, 7
        CALL    EAX2BTR
    CALL NEW_ROW
    MOV     EAX, TESTV
        ADD     EAX, 8
        CALL    EAXCMOV
    CALL NEW_ROW
    MOV     EAX, TESTV
        ADD     EAX, 9
        CALL    EAX2WIN
    CALL NEW_ROW
    MOV     EAX, TESTV
        ADD     EAX, 10
        CALL    EAX2BCD
    CALL NEW_ROW
    RET
TESTR  ENDP

```

```

;-----
MAINN  PROC
    MOV     TESTV, 1000000000
    LL0:
        ;MOV     EDX, TESTV
        ;CALL    EDX2DEC
        MOV     EAX, TESTV
        ;CALL    EAX2ASC
        ;CALL    EAX2DEC
        ;CALL    EAX2AST
        ;CALL    EAX2AFL
        ;CALL    EAX2BTR
        ;CALL    EAXCMOV
        CALL    EAX2BCD
        CALL    S_EQU
        MOV     EAX, TESTV
        CALL    EAX_SQR
        PUSH    EDX
        ;CALL    EAX2DEC
        ;CALL    EAX2BTR
        ;CALL    EAXCMOV
        CALL    EAX2ASC
        ;CALL    EAX2AST
        ;CALL    EAX2AFL
        ;MOV     ESI, EAX
        ;CALL    ESI2ASC
        CALL    S_PLUS
        ;POP     EAX
        ;CALL    EAX2SUB
        POP     EAX
        ;CALL    EAX2ASC
        CALL    EAX2DEC
        ;CALL    EAX2ASF
        ;CALL    EAX2AST
        ;CALL    EAX2AFL
        CALL    NEW_ROW
        INC     TESTV
        CMP     TESTV, 1000000100
        JLE     LL0
    JLE     RET
MAINN  ENDP

```

```

;-----
SPCL   DB      9, 0
TSPACE PROC
    INVOKE StdOut, ADDR SPCL
    RET
TSPACE ENDP

```

```

;-----
NEWROWL DB      13, 10, 0
NEW_ROW PROC
    INVOKE StdOut, ADDR NEWROWL
    RET
NEW_ROW ENDP

```

```

;-----
EPE    DB      253, " + ", 0
S_PLUS PROC
    INVOKE StdOut, ADDR EPE
    RET
S_PLUS ENDP

```

```

;-----
EQE    DB      " = ", 0

S_EQU  PROC
        INVOKE StdOut, ADDR EQE
        RET
S_EQU  ENDP
;-----

D00    DB      02, 1
D01    DB      02, 2
D02    DB      02, 4
D03    DB      03, 8, 0
D04    DB      03, 6, 1
D05    DB      03, 2, 3
D06    DB      04, 4, 6, 0
D07    DB      04, 8, 2, 1
D08    DB      04, 6, 5, 2
D09    DB      05, 2, 1, 5, 0
D10    DB      05, 4, 2, 0, 1
D11    DB      05, 8, 4, 0, 2
D12    DB      05, 6, 9, 0, 4
D13    DB      06, 2, 9, 1, 8, 0
D14    DB      06, 4, 8, 3, 6, 1
D15    DB      06, 8, 6, 7, 2, 3
D16    DB      07, 6, 3, 5, 5, 6, 0
D17    DB      07, 2, 7, 0, 1, 3, 1
D18    DB      07, 4, 4, 1, 2, 6, 2
D19    DB      08, 8, 8, 2, 4, 2, 5, 0
D20    DB      08, 6, 7, 5, 8, 4, 0, 1
D21    DB      08, 2, 5, 1, 7, 9, 0, 2
D22    DB      08, 4, 0, 3, 4, 9, 1, 4
D23    DB      09, 8, 0, 6, 8, 8, 3, 8, 0
D24    DB      09, 6, 1, 2, 7, 7, 7, 6, 1
D25    DB      09, 2, 3, 4, 4, 5, 5, 3, 3
D26    DB      10, 4, 6, 8, 8, 0, 1, 7, 6, 0
D27    DB      10, 8, 2, 7, 7, 1, 2, 4, 3, 1
D28    DB      10, 6, 5, 4, 5, 3, 4, 8, 6, 2
D29    DB      11, 2, 1, 9, 0, 7, 8, 6, 3, 5, 0
D30    DB      11, 4, 2, 8, 1, 4, 7, 3, 7, 0, 1
D31    DB      11, 8, 4, 6, 3, 8, 4, 7, 4, 1, 2

EDX2DEC PROC
        STD
        MOV     DWORD PTR INTRES, 0
        MOV     DWORD PTR INTRES + 4, 0
        MOV     WORD PTR INTRES + 8, 0
        LEA     EBX, D00
        LEA     EDI, 9 + INTRES
        PUT_I_DO:
            MOVZX ECX, BYTE PTR [EBX]
            SHR     EDX, 1
            JNC     PUT_I_END_IF
            INC     EBX
            DEC     ECX
            XOR     AX, AX
            LEA     EDI, 9 + INTRES
            PUT_I_DO1:
                MOVZX AX, AH
                ADD     AL, BYTE PTR [EDI]
                ADD     AL, BYTE PTR [EBX]
                AAA
                STOSB
                INC     EBX
            LOOPD   PUT_I_DO1
            PUT_I_END_IF:
                ADD     EBX, ECX
                TEST    EDX, EDX
            JNZ     PUT_I_DO
            INC     EDI
            OR     DWORD PTR INTRES, "0000"
            OR     DWORD PTR INTRES + 4, "0000"
            OR     WORD PTR INTRES + 8, "00"
            XOR     EDX, EDX
            CMP     BYTE PTR [EDI], "0"
            SETZ    DL
            ADD     EDX, EDI
            CMP     EDX, OFFSET RESULT
            CL
            SETZ    CL
            SUB     EDX, ECX
            CLD
            INVOKE StdOut, EDX
            RET
EDX2DEC ENDP
;-----

ADDIT  DW      "00", "01", "02", "03", "04", "05", "06", "07", "08"
        DW      "09", "10", "11", "12", "13", "14", "15", "16", "17"
        DW      "18", "19", "20", "21", "22", "23", "24", "25", "26"
        DW      "27", "28", "29", "30"
WC     DB      "1"
        DB      02, "2"
        DB      02, "4"
        DB      03, "8", "0"
        DB      03, "6", "1"
        DB      03, "2", "3"
        DB      04, "4", "6", "0"
        DB      04, "8", "2", "1"
        DB      04, "6", "5", "2"
        DB      05, "2", "1", "5", "0"
        DB      05, "4", "2", "0", "1"
        DB      05, "8", "4", "0", "2"
        DB      05, "6", "9", "0", "4"

```

;The routine prints ASCII
 ;decimal content of EDX register
 ;via BCD arithmetic AAA
 ;instruction.
 ;Author: Andrija Radovic, ©2011


```

DB 06, "2", "9", "1", "8", "0"
DB 06, "4", "8", "3", "6", "1"
DB 06, "8", "6", "7", "2", "3"
DB 07, "6", "3", "5", "5", "6", "0"
DB 07, "2", "7", "0", "1", "3", "1"
DB 07, "4", "4", "1", "2", "6", "2"
DB 08, "8", "8", "2", "4", "2", "5", "0"
DB 08, "6", "7", "5", "8", "4", "0", "1"
DB 08, "2", "5", "1", "7", "9", "0", "2"
DB 08, "4", "0", "3", "4", "9", "1", "4"
DB 09, "8", "0", "6", "8", "8", "3", "8", "0"
DB 09, "6", "1", "2", "7", "7", "7", "6", "1"
DB 09, "2", "3", "4", "4", "5", "5", "3", "3"
DB 10, "4", "6", "8", "8", "0", "1", "7", "6", "0"
DB 10, "8", "2", "7", "7", "1", "2", "4", "3", "1"
DB 10, "6", "5", "4", "5", "3", "4", "8", "6", "2"
DB 11, "2", "1", "9", "0", "7", "8", "6", "3", "5", "0"
DB 11, "4", "2", "8", "1", "4", "7", "3", "7", "0", "1"
DB 11, "8", "4", "6", "3", "8", "4", "7", "4", "1", "2"

EAX2BCD PROC
MOV     DWORD PTR INTRES, "0000"
MOV     DWORD PTR INTRES + 4, "0000"
MOV     WORD PTR INTRES + 8, "00"
LEA     EBX, WC
LEA     EDI, 9 + INTRES
EAX2BCD_DO:
MOVZX   ECX, BYTE PTR [EBX]
SHR     EAX, 1
JNC     EAX2BCD_END_IF
INC     EBX
DEC     ECX
MOV     DH, "0"
LEA     EDI, 9 + INTRES
SUB     EDI, ECX
EAX2BCD_D01:
MOVZX   EDX, DH
ADD     DL, BYTE PTR [EDI + ECX]
ADD     DL, BYTE PTR [EBX]
MOV     DX, WORD PTR [2 * EDX + OFFSET ADDIT - 6 * "0"]
MOV     BYTE PTR [EDI + ECX], DL
INC     EBX
LOOPD   EAX2BCD_D01
EAX2BCD_END_IF:
ADD     EBX, ECX
TEST    EAX, EAX
JNZ     EAX2BCD_DO
INC     EDI
XOR     EDX, EDX
CMP     BYTE PTR [EDI], "0"
SETZ    DL
ADD     EDX, EDI
CMP     EDX, OFFSET RESULT
SETZ    CL
SUB     EDX, ECX
INVOKE  StdOut, EDX
RET
EAX2BCD ENDP

;-----
BASES DD 1000000000, 100000000, 10000000, 1000000, 100000, 10000, 1000, 100, 10, 1
NUMDIG DB 9, 9, 9, 8, 8, 8, 7, 7, 7, 6, 6, 6, 6, 5, 5, 5, 4, 4, 4, 3, 3, 3, 3, 2, 2, 2
DB 1, 1, 1, 0, 0, 0

EAX2SUB PROC
MOV     DX, "00" - 101H
BSR     EDI, EAX
CMOVZ   EDI, EAX
MOVZX   EDI, BYTE PTR [EDI + OFFSET NUMDIG]
PUSH    EDI
EAX2SUB_D01:
MOV     EBX, DWORD PTR [4 * EDI + OFFSET BASES]
EBX_D02:
INC     DL
SUB     EAX, EBX
JNC     EBX_D02
MOV     BYTE PTR [EDI + OFFSET INTRES], DL
INC     EDI
MOV     DL, DH
ADD     EAX, EBX
JNZ     EAX2SUB_D01
POP     EDI
CMP     BYTE PTR [EDI + OFFSET INTRES], "0"
SETZ    AL
LEA     EDI, [EDI + EAX + OFFSET INTRES]
CMP     EDI, OFFSET INTRES + 10
SETZ    AL
SUB     EDI, EAX
INVOKE  StdOut, EDI
RET
EAX2SUB ENDP

;-----
ESI2ASC PROC
LEA     EDI, INTRES
MOV     AX, "00" - 101H
PUSH    DWORD PTR 0
PUSH    DWORD PTR 1
PUSH    DWORD PTR 10
PUSH    DWORD PTR 100
PUSH    DWORD PTR 1000
PUSH    DWORD PTR 10000
PUSH    DWORD PTR 100000
PUSH    DWORD PTR 1000000
PUSH    DWORD PTR 10000000
PUSH    DWORD PTR 100000000

```

;The routine prints ASCII
 ;decimal content of EDX register
 ;via BCD arithmetics on the
 ;ADDIT array.
 ;No specific instruction is used.
 ;Author: Andrija Radovic, ©2011

;The routine prints ASCII
 ;decimal content of EAX register
 ;by repetitive subtractions
 ;with coefficients pulled from
 ;array with reduced number
 ;of iteration.
 ;Author: Andrija Radovic, ©2011

;The routine prints ASCII
 ;decimal content of ESI register
 ;by repetitive subtractions
 ;with coefficients pulled from
 ;stack without counter.
 ;Author: Andrija Radovic, ©2011

```

PUSH    DWORD PTR 100000000
MOV     EBX, 1000000000
CLD
EAS_D01:
    EAS_D02:
        INC     AX
        SUB     ESI, EBX
        JNC     EAS_D02
        ADD     ESI, EBX
        STOSB
        MOV     AL, AH
        POP     EBX
        TEST    EBX, EBX
        JNZ     EAS_D01
        MOV     AL, "0"
        MOV     ECX, 10
        LEA     EDI, INTRES
        REPE    SCASB
        DEC     EDI
        INVOKE  StdOut, EDI
        RET
ESI2ASC ENDP
;-----
ASCII_TABLE DW "00", "10", "20", "30", "40", "50", "60", "70", "80", "90", "00", "00"
            DW "00", "00", "00", "00", "01", "11", "21", "31", "41", "51", "61", "71"
            DW "81", "91", "00", "00", "00", "00", "00", "00", "00", "02", "12", "22", "32"
            DW "42", "52", "62", "72", "82", "92", "00", "00", "00", "00", "00", "00", "00"
            DW "03", "13", "23", "33", "43", "53", "63", "73", "83", "93", "00", "00", "00"
            DW "00", "00", "00", "00", "04", "14", "24", "34", "44", "54", "64", "74", "84"
            DW "94", "00", "00", "00", "00", "00", "00", "00", "05", "15", "25", "35", "45"
            DW "55", "65", "75", "85", "95", "00", "00", "00", "00", "00", "00", "00", "06"
            DW "16", "26", "36", "46", "56", "66", "76", "86", "96", "00", "00", "00", "00"
            DW "00", "00", "00", "00", "07", "17", "27", "37", "47", "57", "67", "77", "87"
            DW "97", "00", "00", "00", "00", "00", "00", "00", "08", "18", "28", "38", "48"
            DW "58", "68", "78", "88", "98", "00", "00", "00", "00", "00", "00", "00", "09"
            DW "19", "29", "39", "49", "59", "69", "79", "89", "99", "00", "00"
            DT 0
BCRESUL

EAX2AST PROC
MOV     DWORD PTR BCRESUL, EAX
MOV     DWORD PTR BCRESUL + 4, 0
FILD    QWORD PTR BCRESUL
FBSTP   BCRESUL
LEA     EDI, INTRES
MOV     ECX, 5
CLD
EAX2AST_DO:
MOVZX   EAX, BYTE PTR [ECX + OFFSET BCRESUL - 1]
MOV     AX, WORD PTR [2 * EAX + OFFSET ASCII_TABLE]
        STOSW
        LOOPD  EAX2AST_DO
        MOV     AL, "0"
        MOV     ECX, 10
        LEA     EDI, INTRES
        REPE    SCASB
        DEC     EDI
        INVOKE  StdOut, EDI
        RET
EAX2AST ENDP
;-----
BCDEX    DT 0
EAX2AFL PROC
STD
MOV     DWORD PTR BCDEX, EAX
MOV     DWORD PTR BCDEX + 4, 0
FILD    QWORD PTR BCDEX
FBSTP   BCDEX
MOV     EDX, DWORD PTR BCDEX
MOVZX   EBX, WORD PTR BCDEX + 4
LEA     EDI, 9 + INTRES
EAX2AFL_DO:
MOV     EAX, EDX
AND     EAX, 15
OR      AL, "0"
        STOSB
        SHRD   EDX, EBX, 4
        SHR    EBX, 4
        MOV    EAX, EDX
        OR     EAX, EBX
        JNZ    EAX2AFL_DO
        LEA    EDX, [EDI + 1]
        CLD
        INVOKE StdOut, EDX
        RET
EAX2AFL ENDP
;-----
EAX2DEC PROC
LEA     ECX, INTRES + 9
MOV     EBX, 10
EAX2DEC_DO:
XOR     EDX, EDX
DIV     EBX
OR      EDX, "0"
        MOV    BYTE PTR [ECX], DL
        TEST   EAX, EAX
        LOOPNZ EAX2DEC_DO
        LEA    EDX, [ECX + 1]
        INVOKE StdOut, EDX
        RET
EAX2DEC ENDP
;-----
;The routine prints ASCII
;decimal content of EAX register
;by usage of coprocessor FBSTP
;instruction which converts
;number into the packed BCD.
;Unpacking of BCD is done via
;the table of unpacked values.
;Author: Andrija Radovic, @2011

;The routine prints ASCII
;decimal content of EAX register
;by usage of coprocessor FBSTP
;instruction which converts
;number into the packed BCD.
;Unpacking of BCD is done via
;SHR and SHRD instructions.
;Author: Andrija Radovic, @2011

;This routine prints ASCII
;decimal content of EAX register
;by its repetitive dividing by
;10 using DIV instruction that
;simultaneously yields result
;and remainder which denotes
;current decimal digit.
;Author: Andrija Radovic, @2011

```

```

;
EAX2ASC PROC                                     ;This routine prints ASCII
LEA      ECX, INTRES + 9                        ;decimal content of EAX register
MOV      EDI, 858993459                        ;by its repetitive dividing by
MOV      EBX, EAX                              ;10 using MUL instruction to
AX2ASC_D0:                                     ;divide by 10 via multiplication
LEA      EAX, [EBX + 1]                        ;with the appropriate constant.
MUL      EDI                                  ;Author: Andrija Radovic, @2011
SHR      EDX, 1
LEA      EAX, [4 * EDX + EDX]
NEG      EAX
LEA      EAX, [EBX + 2 * EAX + "0"]
MOV      BYTE PTR [ECX], AL
MOV      EBX, EDX
LOOPNZD AX2ASC_D0
LEA      EDX, [ECX + 1]
INVOKE   StdOut, EDX
RET
EAX2ASC ENDP

;-----
LB DB 0F0H,0F1H, 1,3, 0F2H,4, 0F3H,0F4H, 2,7, 0F5H,0F6H, 6,8, 0F7H,9, 0F8H,0F9H
D DD 0,1,2,3,4,5,6,7,8,9
DD 00,10,20,30,40,50,60,70,80,90
DD 000,100,200,300,400,500,600,700,800,900
DD 0000,1000,2000,3000,4000,5000,6000,7000,8000,9000
DD 00000,10000,20000,30000,40000,50000,60000,70000,80000,90000
DD 000000,100000,200000,300000,400000,500000,600000,700000,800000,900000
DD 0000000,1000000,2000000,3000000,4000000,5000000,6000000,7000000,8000000,9000000
DD 00000000,10000000,20000000,30000000,40000000,50000000,60000000,70000000,80000000,90000000
DD 000000000,100000000,200000000,300000000,400000000,500000000,600000000,700000000,800000000,900000000
DD 0000000000,1000000000,2000000000,3000000000,4000000000,5000000000,6000000000,7000000000,8000000000,9000000000
DD 00000000000,10000000000,20000000000,30000000000,40000000000,50000000000,60000000000,70000000000,80000000000,90000000000
DD 000000000000,100000000000,200000000000,300000000000,400000000000,500000000000,600000000000,700000000000,800000000000,900000000000
DD 0000000000000,1000000000000,2000000000000,3000000000000,4000000000000,5000000000000,6000000000000,7000000000000,8000000000000,9000000000000
DD 00000000000000,10000000000000,20000000000000,30000000000000,40000000000000,50000000000000,60000000000000,70000000000000,80000000000000,90000000000000
DD 000000000000000,100000000000000,200000000000000,300000000000000,400000000000000,500000000000000,600000000000000,700000000000000,800000000000000,900000000000000
DD 0000000000000000,1000000000000000,2000000000000000,3000000000000000,4000000000000000,5000000000000000,6000000000000000,7000000000000000,8000000000000000,9000000000000000
DD 00000000000000000,10000000000000000,20000000000000000,30000000000000000,40000000000000000,50000000000000000,60000000000000000,70000000000000000,80000000000000000,90000000000000000
DD 000000000000000000,100000000000000000,200000000000000000,300000000000000000,400000000000000000,500000000000000000,600000000000000000,700000000000000000,800000000000000000,900000000000000000
DD 0000000000000000000,1000000000000000000,2000000000000000000,3000000000000000000,4000000000000000000,5000000000000000000,6000000000000000000,7000000000000000000,8000000000000000000,9000000000000000000
DD 00000000000000000000,10000000000000000000,20000000000000000000,30000000000000000000,40000000000000000000,50000000000000000000,60000000000000000000,70000000000000000000,80000000000000000000,90000000000000000000
DD 000000000000000000000,100000000000000000000,200000000000000000000,300000000000000000000,400000000000000000000,500000000000000000000,600000000000000000000,700000000000000000000,800000000000000000000,900000000000000000000
DD 0000000000000000000000,1000000000000000000000,2000000000000000000000,3000000000000000000000,4000000000000000000000,5000000000000000000000,6000000000000000000000,7000000000000000000000,8000000000000000000000,9000000000000000000000
DD 0000000000000000000000,10000000000000000000000,20000000000000000000000,30000000000000000000000,40000000000000000000000,50000000000000000000000,60000000000000000000000,70000000000000000000000,80000000000000000000000,90000000000000000000000
DD 00000000000000000000000,100000000000000000000000,200000000000000000000000,300000000000000000000000,400000000000000000000000,500000000000000000000000,600000000000000000000000,700000000000000000000000,800000000000000000000000,900000000000000000000000
DD 000000000000000000000000,1000000000000000000000000,2000000000000000000000000,3000000000000000000000000,4000000000000000000000000,5000000000000000000000000,6000000000000000000000000,7000000000000000000000000,8000000000000000000000000,9000000000000000000000000
DD 0000000000000000000000000,10000000000000000000000000,20000000000000000000000000,30000000000000000000000000,40000000000000000000000000,50000000000000000000000000,60000000000000000000000000,70000000000000000000000000,80000000
```

```

JNS     TEST     AX, AX
EAXCMOV_D01
MOVZX   EAX, AL
LEA     EDI, 9 + INTRES
SUB     EDI, EAX
PUSH    EDI
LEA     ESI, [EAX + 4 * EAX]
MOV     EAX, 5
SHL     ESI, 3
EAXCMOV_D03:
CMP     EBX, DWORD PTR [4 * EAX + ESI + OFFSET DC]
CMOVC   AX, WORD PTR [4 * EAX + (OFFSET LC - 4)]
CMOVNC  AX, WORD PTR [4 * EAX + (OFFSET LC - 2)]
;These two instructions
;work on some processors...

CMOVC   DX, WORD PTR [4 * EAX + (OFFSET LC - 4)]
CMOVNC  DX, WORD PTR [4 * EAX + (OFFSET LC - 2)]
MOV     AX, DX
;Malfunction of x86 CMOV
;requires these ones instead
;of above two instructions.

JNS     TEST     AX, AX
EAXCMOV_D03
MOVZX   EAX, AL
SUB     EBX, DWORD PTR [4 * EAX + ESI + OFFSET DC]
OR      AL, "0"
STOSB
MOV     EAX, 5
SUB     ESI, 40
JNC     EAXCMOV_D03
POP     EDX
INVOKE  StdOut, EDX
RET
EAXCMOV ENDP

;-----

NumFormat DB      "%u", 0
Bufferw   DB      32 DUP(0)

EAX2WIN PROC
INVOKE    wsprintf, ADDR Bufferw, ADDR NumFormat, EAX
INVOKE    StdOut, ADDR Bufferw
RET
EAX2WIN ENDP

;-----

EAX_SQR PROC
XOR     ESI, ESI
XOR     EDX, EDX
MOV     EBX, 1073741824
SQRD0:
LEA     EDI, [EBX + ESI]
ADD     EDI, EDX
SHR     EDX, 1
CMP     EAX, EDI
JC      SQRD_END_IF
MOV     ESI, EDI
ADD     EDX, EBX
SQRD_END_IF:
SHR     EBX, 2
JNZ     SQRD0
SUB     EAX, ESI
XCHG   EAX, EDX
RET
EAX_SQR ENDP

;-----

END      ASCII

```

Above routines should be assembled rather from the command line than from some assembler's IDE because most of IDEs usually do not offer full control of the parameters passed to assembler and linker, just as is required by the linking process with removed section's read-only property.

Therefore they should be assembled manually from the command line:

```

C:\masm32\bin\ML.EXE /c /coff /Cp /nologo
/I"\Masm32\Include" ASCIIwin.asm

```

And it should be linked with the following single line which is spread to two lines of text although this actually should be only one uninterrupted line in the console:

```

C:\masm32\bin\LINK.EXE /SUBSYSTEM:CONSOLE /RELEASE
/LIBPATH:"\Masm32\Lib" /section:".text",ERW ASCIIwin.obj

```

Above two lines make EXE program with writable Code section. The benefit of such hand crafted linking is ability to write self modifying code.

Although it is highly discouraged by the MASM manual, this is something that should not be instantly rejected! Actually, all the recommendations those allegedly improve prettiness of the code by forcing real limitations to programmers should be immediately rejected! The self modifying code is very useful and sometimes may shrink code a lot and rejection of such option just because somebody found that it is not fancy enough seems to be quite wrong way of action. The second ability to keep all the routines' peaces tightly together is also very important for quick and efficient coding, which all requires section with the code to be writable.

CONCLUSION

All presented methods suffer from some sorts of disadvantages that elongate their executions. The misery of the task is consisted of the fact that set of 2 elements should be translated into the set of only 10 elements with up to 10 decimal positions and this exceptionally small number of decimal digits and decimal positions (32-bit integer has 10 decimal positions) really limits the efficiency of all of the applied optimization schemes, highly challenging coding skills because the number of instructions in optimization schemes reach the number of instructions performed in main loops.

Although the algorithm **EDX2DEC** based on the **AAA** instruction altogether with the **EAX2BTR** based on binary tree search are duly incomprehensible at the glance, this give it special charm making them quite fashionable. Despite the fact that **EAX2BTR** utilizes theoretically very promising Binary Tree Search, it is not much better than much simpler algorithms like **EAX2SUB** based on the repetitive subtractions. The one non-trivial algorithm with comparable performance with simplest ones is the **EAX2ASC** routine that utilizes division trough multiplication with only 8 instructions within main loop, which is going to be quite fast on the new processors with **MUL** instruction able to accomplish it in just of a couple of clock cycles. The algorithm **EAX2DEC** based on the **DIV** instruction is absolute champion of shortness – only five instructions within the main loop, but **DIV** instruction may be very slow depending on the particular processor. The situation is quite reversal when these algorithms should be transferred in VHDL hardware, then the binary tree search and BCD **AAA** instruction based algorithms are the most promising ones offering tremendous opportunities for hardware parallelization – they can reach even execution in a single clock cycle only.

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